

CS500
Spring 2018, Assignment #3

PROBLEM 6 (1+3+3+3P) :

Consider the following algorithm, regardless of its purpose:

Given an n -tuple (b_1, \dots, b_n) of bits, scan for the (index j of the) first bit that is non-zero. If all are zero, count to $2^n - 1$. Then stop.

- a) Analyze the worst-case cost of this algorithm, asymptotically as $n \rightarrow \infty$.
- b) Analyze the average-case cost of this algorithm, asymptotically as $n \rightarrow \infty$.
- c) Consider a binary counter, initialized with zero and supporting *two* operations: increment and decrement; $\text{decrement}(0)=0$. Analyze the amortized number of bit-flips. Prove that your asymptotic estimate is tight.
- d) Consider a binary counter initialized with any ℓ -bit number and then incremented n times. Devise a potential function to prove that the amortized number of bit-flips is $\mathcal{O}(1 + \ell/n)$.

PROBLEM 7 (5+5P):

- a) Draw the Fibonacci Heap (with marks) resulting from the following sequence of operations:

```
create(9), insKey(8), insKey(7), extrMin(), insKey(6),  
insKey(5), insKey(4), extrMin(), delKey(6), insKey(3),  
insKey(2), insKey(1), extrMin(), delKey(3), insKey(3).
```

- b) Draw the final disjoint forest (including ranks) resulting from the following sequence of operations according to *Lazy Union-by-Rank** and *Path Compression*:

```
make(1), make(2), make(3), make(4), make(5), ..., make(8), make(9),  
union(1,2),  $x := \text{find}(1)$ , union( $x$ ,3), union( $x$ ,4), union( $x$ ,5), union(6,7),  
union(8,9), union( $\text{find}(6)$ , $\text{find}(8)$ ), union( $\text{find}(7)$ , $x$ ), find(3)
```

*When taking the union of two trees of same rank, let ties be broken s.t. the *second* one becomes *child* of the first.