

Syllabus

4. Intro Data Structures

- Hardware/Math Data Types
- Abstract Data Types
 - Specification
 - Primitives:
semantics and cost
 - Design
 - Analysis
 - (Optimality)
- Basic: Boolean, Integer
- Derived: Array, Queue, Stack
- Linked Data Structures
- (Balanced) Search Trees
- AVL Trees

Hardware vs. Math. Data Types

Martin
Ziegler

- Each country/leader described by "A.I." character parameters.
- Initially *Gandhi*.**aggression** := 1
- When country adopts democracy, **aggression** -= 2.

(signed)
byte/word/dword
 $\neq (\mathbb{N}) \mathbb{Z}$



Recap from Logic / *Model Theory*

Boolean \mathbb{B}	lattice	$0, 1, \wedge, \vee, \neg$
Integer \mathbb{Z}	ring	$0, 1, +, -, \times, \text{div}, >$
Real \mathbb{R}	field	$0, 1, +, -, \times, \div, >$
Complex \mathbb{C}	field	$0, 1, +, -, \times, \div$
polynomials $\mathbb{R}[X]$	ring	$0, 1, +, -, \times, \text{div}$
powerset $\mathcal{P}(S)$	lattice	$\emptyset, S, \cap, \cup, \setminus$
Grassmannian $\mathcal{G}(V)$	lattice	$\{0\}, V, \cap, \oplus, \perp$
<i>Presburger</i> Integers \mathbb{Z}	group	$0, 1, +, -, >$

Meta-Def: A *structure* is a set, together with some

- constants,
- functions, and
- relations

Axioms are propositions satisfied by said structure.

Abstract Data *Types*

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Ziegler

Boolean \mathbb{B}

$0, 1, \wedge, \vee, \neg$

Integer \mathbb{Z} (vs. byte/word etc.)

$0, 1, +, -, \times, \text{div}, >$

Real \mathbb{R} (vs. float/double etc.)

Stack of X

$\text{new}, \text{push } x, \text{pop}, \text{isEmpty}$

Queue of X

$\text{new}, \text{enqueue } x, \text{dequeue}, \text{isEmpty}$

FIFO

LIFO:

$\text{pop} \circ \text{push} = \text{id}$

Definition: An *abstract data type* is a “class of objects” together with certain operations.

Axioms describe their behavior

regardless of
implementation

Data Structures and Cost

Martin
Legler

Boolean \mathbb{B}

$O(1)$

0, 1, \wedge , \vee , \neg

Integer \mathbb{Z} (vs. byte/word etc.)

0, 1, +, -, \times , div, >

Real \mathbb{R} (vs. float/double etc.)

$O(n \cdot \log n)$

Stack of X

$O(n) =$
 $O(\log |x|)$

new, push x , pop, isEmpty

Queue of X

new, enqueue x , dequeue, isEmpty

(Dynamic) array of X

$O(1) ?$

[], size, (re-size), search

$O(n)$

Sorted array of Y

$O(\log n)$

search y , insert y , delete y

Priority queue of Y

findmin, insert y

$O(n)$

where Y is totally ordered

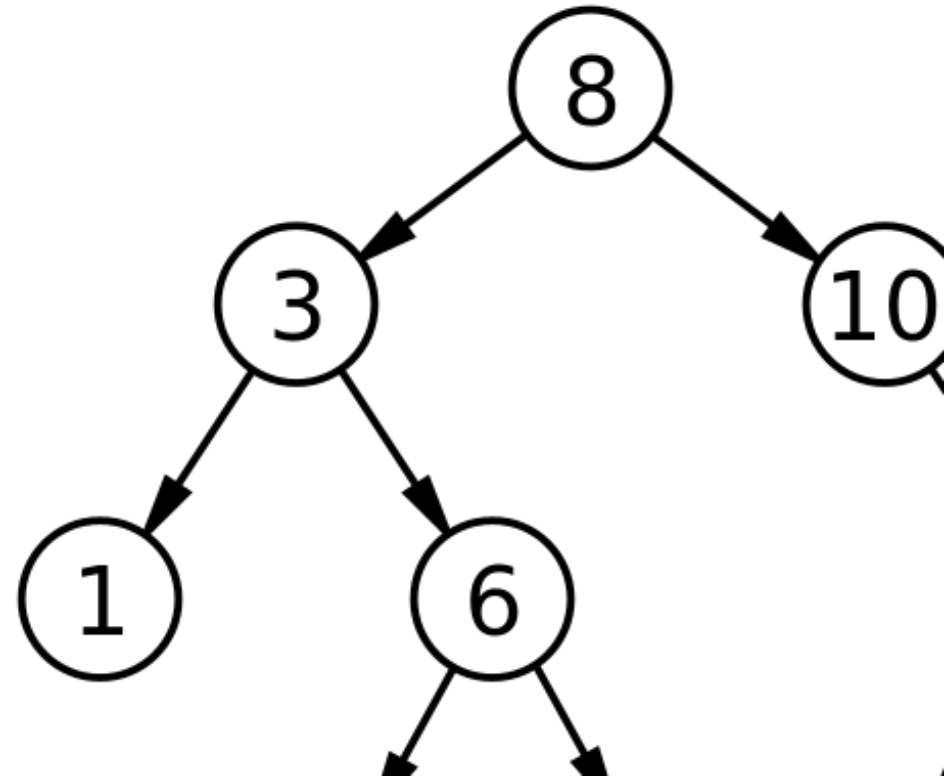
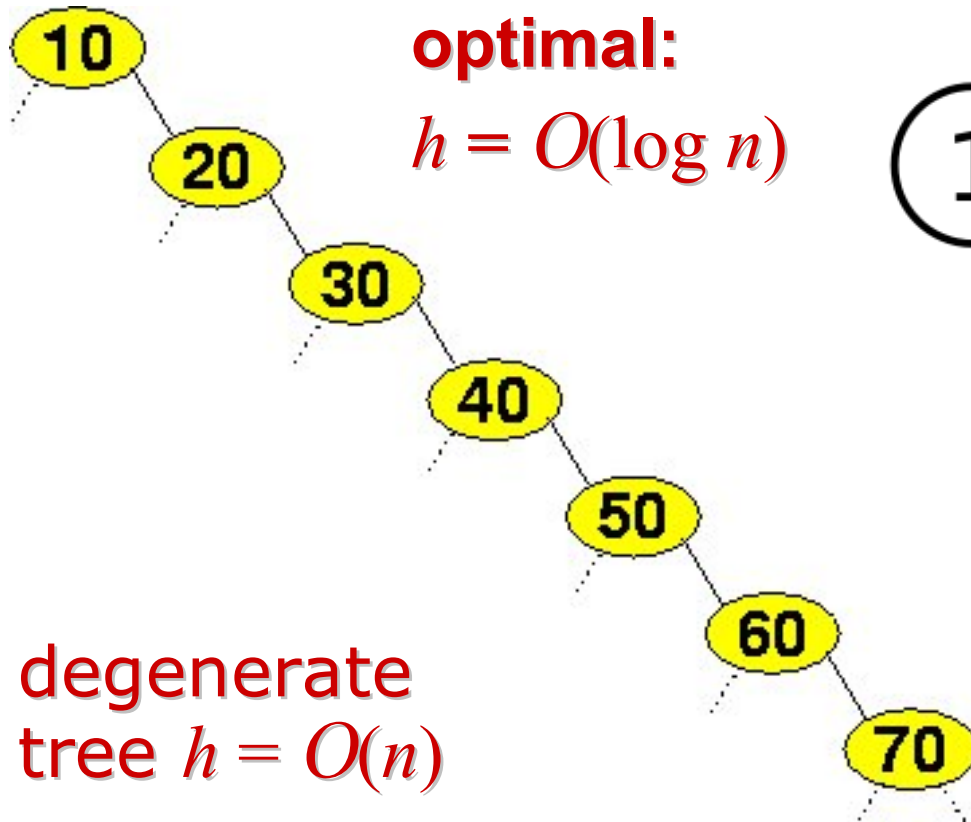
$O(\log n)$

Binary Trees

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Lemma: $n \leq 2^{h+1} - 1$

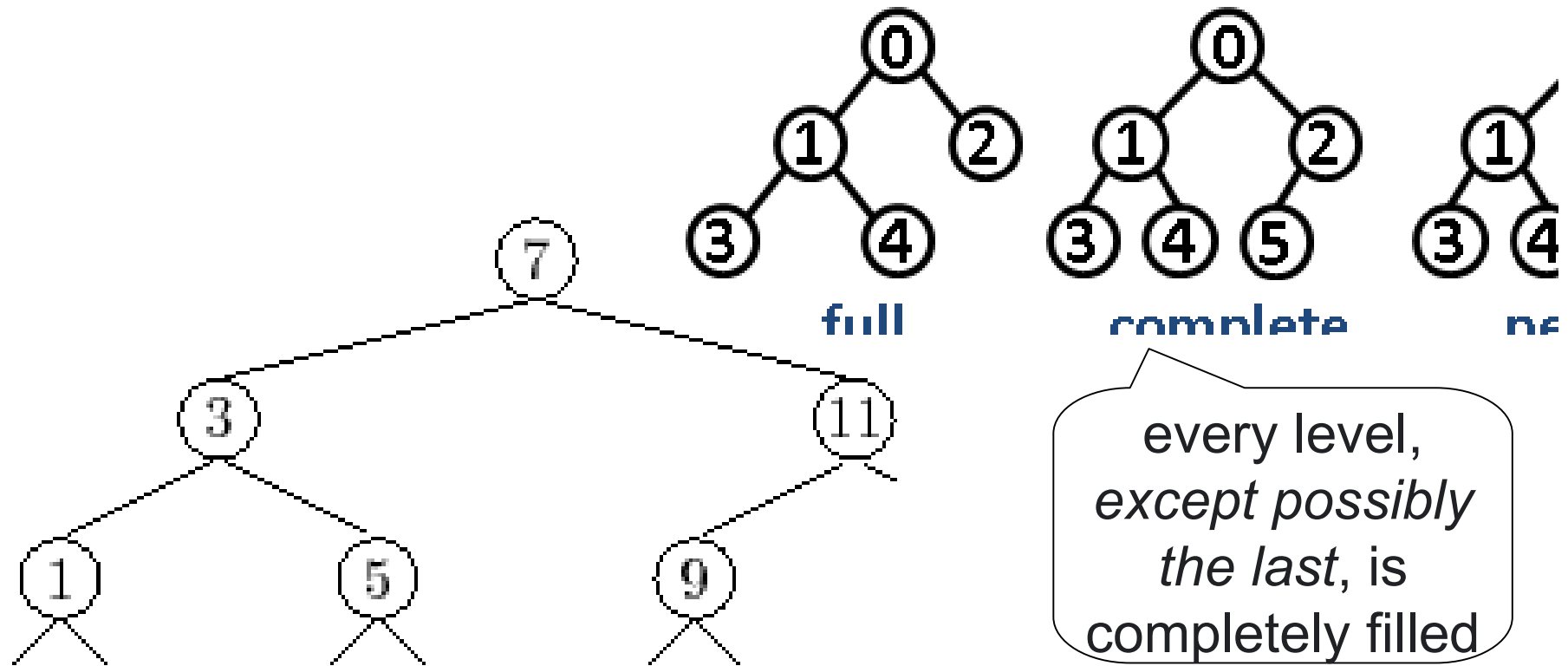
$\Rightarrow h \geq O(\log n)$



Binary search tree:
search, insert, delete
in $O(h)$, $h = \text{height}$

Balanced Binary Trees

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Now **insert(15)**
and **delete(0)**:
Rebalancing is costly!

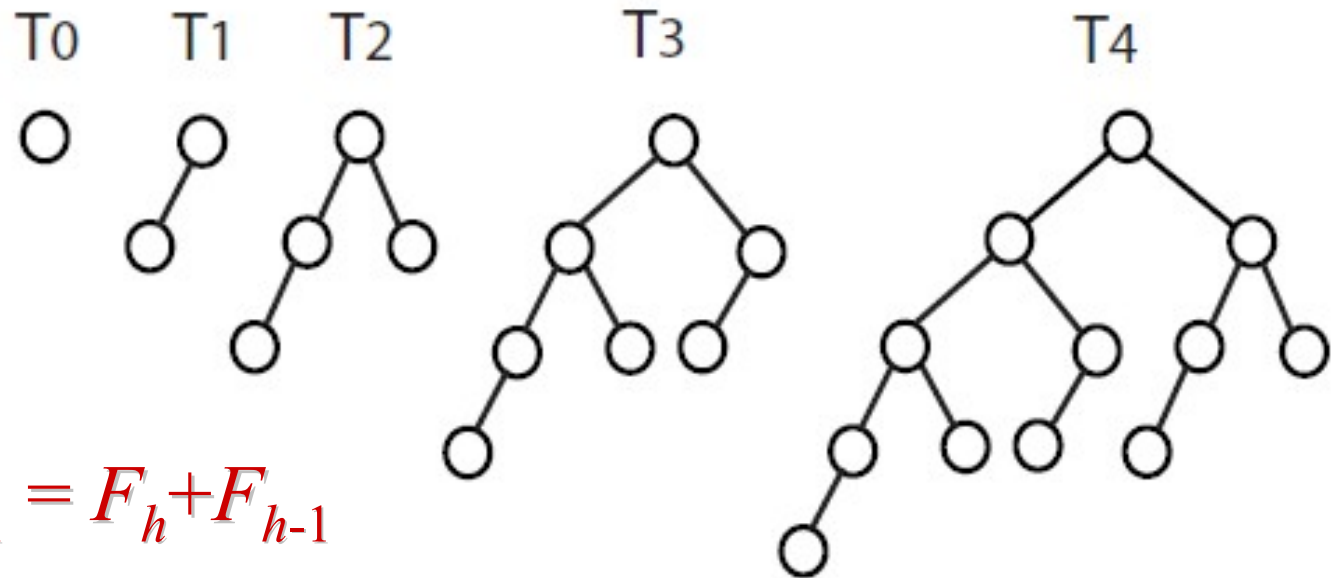
Binary search tree:
search, insert, delete
in $O(h)$, $h = \text{height}$

Adelson-Velsky-Landis'62

Martin
Legler

Binary tree s.t. *left and right subtrees have height difference at most 1!*

minimal
AVLTrees:



Fibonacci

$$F_0=0, F_1=1, F_{h+1} = F_h + F_{h-1}$$

$n(h) := \min$ #nodes of AVLTree of height $h \leq O(\log n)$

$$\#n(0)+1 = F_3, \quad \#n(h+1) + 1 = \#n(h) + 1 + \#n(h-1) + 1 = F_{h+4}$$

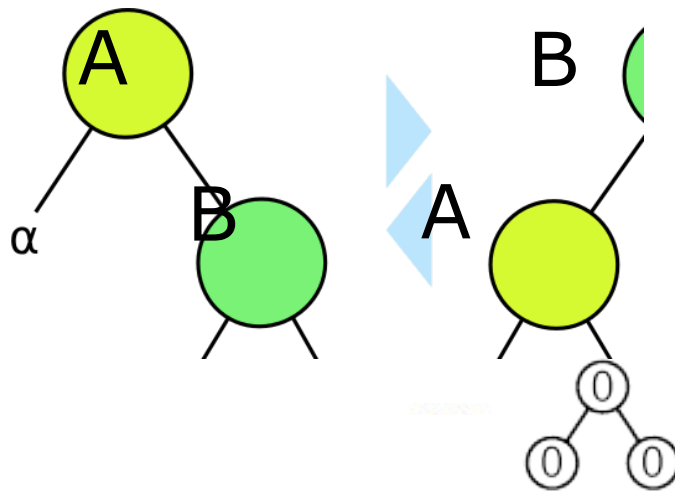
$$\text{Recall } F_h = (\varphi^h - (-1/\varphi)^h) / \sqrt{5} \geq \Omega(1.6^h) \Rightarrow h = O(\log F_h)$$

AVL Tree Maintenance

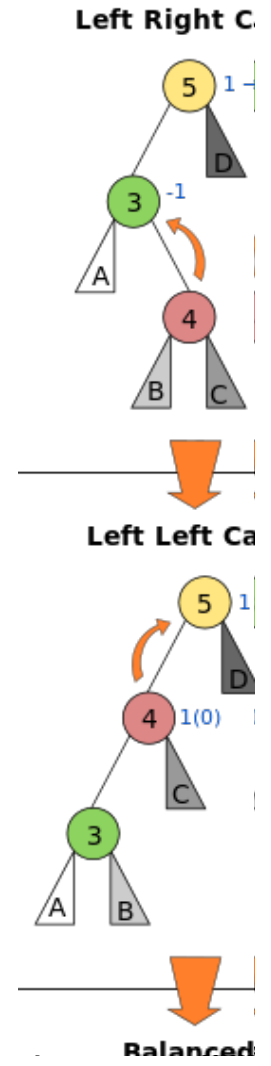
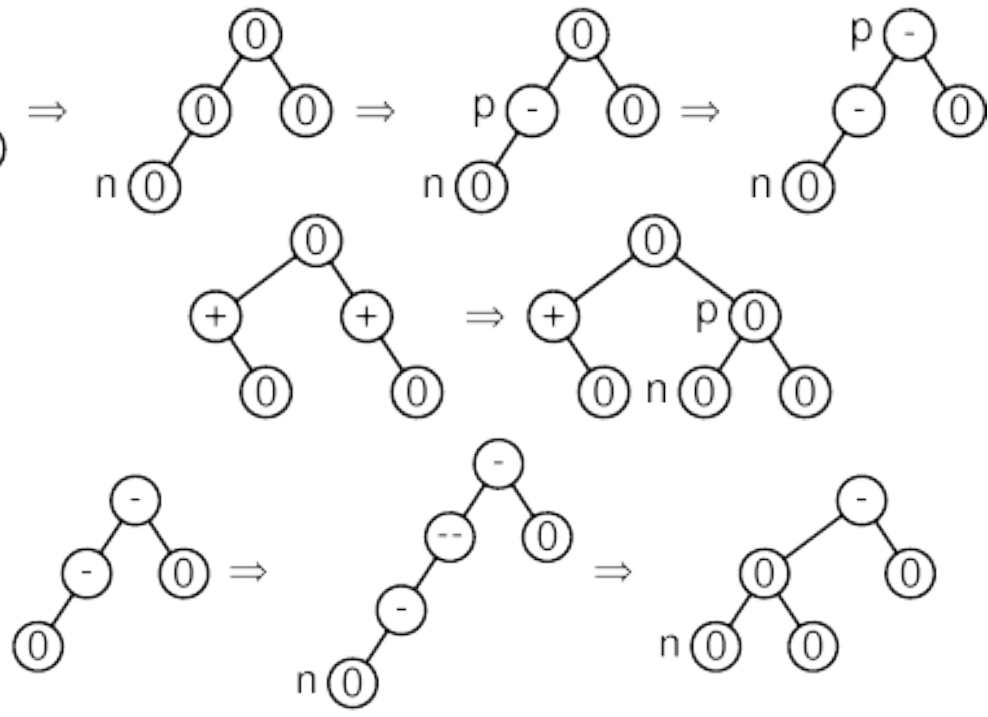
Martin Ziegler

Binary tree s.t. *left and right subtrees have height difference at most 1!*

$$h \leq O(\log n)$$



Store & recursively update balance indicators +, 0, -. After **insert** at a left leaf, propagate up: three cases



search $O(\log n)$
 insert $O(\log n)$
 delete $O(\log n)$
merge $O(n)$

Recap

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