Martin Ziegler Ivan Adrian Koswara Issued on: May 2, 2019 Solutions due: May 13, 2019

CS422

Spring 2019, Homework #4

Problem 10 ($7 \times 3 + 1$ points)

Let *G* be the 3×3 rook graph, pictured on the right. Answer the following questions about *G* and prove your answers.

- a. Is G 3-colorable?
- b. What is the size of the smallest edge cover of *G*?
- c. What is the size of the smallest vertex cover of *G*?
- d. Does *G* have an Eulerian cycle?
- e. Does *G* have a Hamiltonian cycle?
- f. What is the size of the largest clique in *G*?
- g. What is the size of the largest independent set in *G*?
- h. (Bonus) Do you know why it is called a rook graph? (Feel free to look this up.)

Problem 11 (5×3 points)

Consider the following Boolean formula φ :

$$(\neg (a \lor b \lor \neg c) \lor (a \land \neg b))$$

Answer the following questions about φ and prove your answers.

- a. Is φ in CNF? If not, convert it into an equivalent CNF formula.
- b. Is φ in DNF? If not, convert it into an equivalent DNF formula.
- c. Is φ satisfiable?
- d. Is φ a tautology?
- e. Does φ have an equivalent shorter formula? (The length of a formula is the number of times you need to invoke the construction rules. For example, the formula $(a \lor \neg b)$ has length 4: one to make *a*, one to make *b*, one to make $\neg b$, and one to make the entire formula.)

(Problems 12-13 on next page)



1 2 1 2 1 3



Define the decision problem SUDOKU as follows: given (the encoding of) a Sudoku puzzle, does it have a solution?

Define the decision problem SUDOKU_k as follows: given (the encoding of) a $k^2 \times k^2$ Sudoku puzzle, does it have a solution?

	4		
		2	
			2
1			

- a. Does the Sudoku puzzle on the left belong to the problem SUDOKU? Prove it.
- b. Suppose we change the 1 on the bottom-left cell into a 2. Does your answer for item a) change? Prove it.
- c. Prove that SUDOKU is in **NP**.
- d. For any fixed k, prove that SUDOKU_k is in **P**.

Problem 13 (6 + 6 points)

In this problem, you will see that the relationship between **P** and **NP** is very similar to the relationship between decidable and semidecidable problems: we can use a definition for **NP** and adjust it slightly to get an alternative definition for semidecidability.

Consider the set $L_p = \{x | \exists w. \ell(w) \le \text{poly}(\ell(x)) \land \langle x, w \rangle \in L'_p\}$ where $L'_p \in \mathbf{P}$. This is the definition of a problem that is in **NP**: an input is in L_p if has a polynomial-length witness w that can be verified in polynomial time, where L'_p means "is w a correct witness for x?"

Now consider what happens if we get rid of all the polynomials (marked in dark red above): suppose *L* can be expressed as $\{x | \exists w. \langle x, w \rangle \in L'\}$ where *L'* is decidable.

- a. Prove that *L* is semidecidable.
- b. Conversely, prove that any semidecidable problem can be expressed in this form.

Problem 12 (3 + 3 + 6 + 6 points)

In a **Sudoku puzzle**, you are given an $k^2 \times k^2$ grid divided into k^2 boxes of size $k \times k$. Some cells have been filled with integers from the set $\{1, 2, ..., k^2\}$. A **solution** to a Sudoku puzzle is a completely filled $k^2 \times k^2$ grid such that whenever there is a number in a cell of the puzzle, the same number is on the same cell in the solution, and every row, column, and box contains each of $1, 2, ..., k^2$ exactly once. See the pictures on the right; the top one is a 4×4 Sudoku puzzle (k = 2) and the bottom one is its solution. You can verify that the solution is correct.